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(54) Handshaking circuit for resolving contention on a transmission medium regardless of its length

(57) For establishing a parent-child relationship between adjacent network nodes interconnected by a transmission medium, a handshaking circuit comprises a state detector for monitoring the transmission medium to detect predetermined first and second states and a contention that occurs when predetermined first states are simultaneously present. A state machine asserts the first state for initiating a handshaking process and relinquishes this state when a contention is detected. A receive count value is continuously incremented from the instant the contention is detected to the instant the transmission medium changes to an idle state. A transmit count value is continuously incremented from the instant the first state is detected to the instant the contention is detected. When the transmit count value is greater than the receive count value, the state machine asserts the first state again and when the transmit count value is smaller than the receive count value, it asserts the second state.

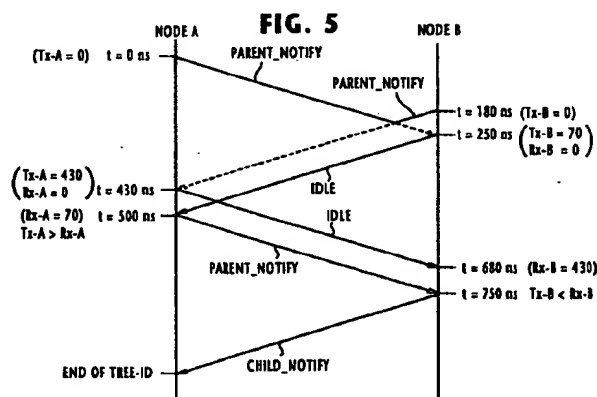
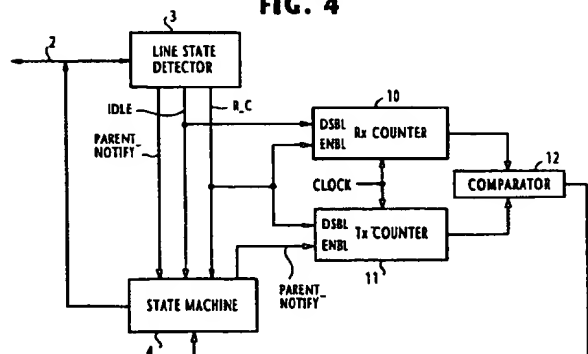


FIG. 4



## Description

The present invention relates to a handshaking circuit for a network of interconnected nodes for establishing a particular relationship, such as parent-child relationship, between adjacent nodes by resolving a contention therebetween which occurs on a transmission medium, such as serial bus, when they simultaneously assert the same protocol status.

A handshaking protocol is standardized by the IEEE Computer Society for interconnecting nodes such as personal computers and associated electronic devices through their ports using a serial bus known as IEEE 1394. Adjacent nodes first establish a parent-child relationship by having one of the nodes assert itself as a "child" of the other. In response, the other node recognizes itself as the parent of the child node and sends an acknowledgment. However, if the parent node has already asserted itself as a child node immediately before it recognizes the parenthood, a contention (collision) is the to occur between them. When both nodes detect this contention, they relinquish their assertion and wait a random period before making a second attempt.

However, the maximum random period is not sufficient to cover a bus length longer than the standard length of 4.5 meters. If the length of a serial bus is much longer than the standard length, there are instances where the contention cannot be resolved.

One solution to this problem would be to increase the time difference between the short and long random period values. Since the node is held in an inactive state during the time its random period timer is still counting clock, the use of a large fixed random value would result in an unnecessarily long interval to resolve a contention on short-length serial buses.

It is therefore an object of the present invention to provide a handshaking circuit capable of resolving an inter-node contention regardless of the length of the transmission medium.

According to a first aspect of the present invention, there is provided a handshaking circuit for establishing a particular relationship between nodes interconnected by a transmission medium. The handshaking circuit comprises a state detector for monitoring the transmission medium to detect a predetermined first state, a predetermined second state, a contention when predetermined first states are simultaneously asserted on the transmission medium, and an idle state when either of the first and second states is not asserted on the transmission medium. A state machine is connected to the state detector for asserting the first state on the transmission medium to initiate a handshaking process and relinquishing the first state when the contention is detected by the state detector. A receive counter is connected to the state detector for continuously incrementing a receive count value from the instant the contention is detected to the instant the idle state is detected. A

transmit counter is connected to the state detector and the state machine for continuously incrementing a transmit count value from the instant the first state is asserted to the instant the contention is detected. The receive count value is compared with the transmit count value. When the transmit count value is greater than the receive count value, the state machine asserts the first state on the transmission medium again and when the transmit count value is smaller than the receive count value, it asserts the second state on the transmission medium.

According to a second aspect, the present invention provides a handshaking circuit for establishing a particular relationship between nodes interconnected by a transmission medium. The handshaking circuit comprises a state detector for monitoring the transmission medium to detect a predetermined first state, a predetermined second state, a contention when predetermined first states are simultaneously asserted on the transmission medium and an idle state when either of the first and second states is not asserted on the transmission medium. A state machine is connected to the state detector for asserting the first state on the transmission medium to initiate a handshaking process and relinquishing the first state when the contention is detected by the state detector. A receive counter is connected to the state detector for continuously incrementing a receive count value from the instant the contention is detected to the instant the idle state is detected. A transmit counter is connected to the state detector and the state machine for continuously incrementing a transmit count value from the instant the first state is asserted to the instant the contention is detected. A decision circuit produces a first decision output when the transmit and receive count values are lower than a predefined threshold and a second decision output when one of the transmit and receive count values is greater than the predefined threshold. A random period timer is activated in response to the detection of the contention by the state detector. The state machine is arranged to respond to the first decision output for asserting the first state on the transmission medium if the random period timer expires during the absence of the first and second states on the transmission medium and asserting the second state on the transmission medium if the random period timer expires during the presence of the first state on the transmission medium. The state machine responds to the second decision output for asserting the first state on the transmission medium when the transmit count value is greater than the receive count value and asserting the second state on the transmission medium when the transmit count value is smaller than the receive count value.

The present invention will be described in further detail with reference to the accompanying drawings, in which:

Fig. 1 is a block diagram of a network of nodes

interconnected by serial buses according to the IEEE standard 1394;

Fig. 2 is a block diagram of a prior art handshaking circuit;

Fig. 3A is a time sequence diagram of the current procedure when the length of an internode transmission medium is within the specified maximum value;

Fig. 3B is a time sequence diagram of the current procedure when the length of an internode transmission medium exceeds the specified maximum value;

Fig. 4 is a block diagram of a handshaking circuit according to one embodiment of the present invention;

Fig. 5 is a time sequence diagram of the handshaking circuits of Fig. 4 during a tree identify process between two contending nodes;

Fig. 6 is a block diagram of a handshaking circuit according to a further embodiment of the present invention;

Fig. 7 is a time sequence diagram of the handshaking circuits of Fig. 6 during a tree identify process between two contending nodes;

Fig. 8 is a block diagram of a handshaking circuit according to a modified embodiment of the present invention;

Fig. 9 is a time sequence diagram of the handshaking circuits of Fig. 8 during a tree identify process between two contending nodes;

Fig. 10 is a schematic diagram of a network using repeaters between two nodes according to a further embodiment of the present invention;

Fig. 11 is a block diagram of each repeater of Fig. 10;

Fig. 12 is a block diagram of each handshaking circuit of Fig. 10; and

Fig. 13 is a time sequence diagram of a tree identify process of the repeaters of Fig. 11 and the prior art handshaking circuits of Fig. 12.

Referring to Fig. 1, there is shown a network of nodes 1-1 to 1-5, such as personal computers and associated electronic devices, which are interconnected by serial buses 2 on which signals of the format standardized by the IEEE Computer Society under "IEEE Std 1394-1995", (or simply "the IEEE 1394") are transmitted. Each of the network nodes uses one or more ports for daisy-chain connection or branch connection to other nodes. Nodes 1-2, 1-4 and 1-5 use only one port for connection to another node and therefore they are called "leaf nodes" and nodes 1-1 and 1-3 use more than one port for connection to other nodes and therefore they are called "branch nodes".

The IEEE 1394 specifies a procedure for determining the parent-child relationships between adjacent nodes. This procedure, known as the tree identify process, begins following a bus reset procedure to configure

the network into a tree with a unique root node and labels such port as connected to either the parent of a node or its children.

Before proceeding with the detailed description of the present invention, it may prove helpful to provide an explanation of the prior art tree identify process with reference to Figs. 2, 3A and 3B.

Each node has a handshaking circuit as shown in Fig. 2. The handshaking circuit includes a line state detector 3 and a notification circuit, or state machine 4, which connected to the serial bus, or cable 2. The line state detector 3 monitors incoming signals on the bus 2 and produces state indication signals. The state machine algorithm of the tree identify process is embodied in the state machine 4 to transmit a handshaking signal onto the output line of cable 2 by holding the state of the cable at one of Parent\_Notify, Child\_Notify and idle (i.e., absence of signal) according to the protocol of the tree identify process. When the state machine 4 changes the line state to Parent\_Notify following the detection of Parent\_Notify by the line state detector 3, a contention, known as a Root\_Contention, is said to exist on the cable between adjacent nodes.

When a root contention occurs, the line state detector 3 produces a contention presence signal. In response, the state machine 4 relinquishes the Parent\_Notify state and changes the line state to idle, and a counter 6 starts incrementing its clock count and supplies its output to a comparator 7. A random number generator 9 is activated in response to the contention presence signal and supplies a random number, which is binary 0 or 1, to a threshold generator 8. Threshold generator 8 assumes one of two threshold values depending on the random number. If the random number is 0, the threshold generator 8 produces a lower timeout threshold corresponding to the period of 260 ns (which is specified as Contend\_Fast according to the tree identify protocol) and if the random number is 1, it produces a higher timeout threshold corresponding to the period of 600 ns (specified as Contend\_Slow). The comparator 7 compares the output of the counter 6 with the random threshold value to produces an output when the clock count exceeds the random threshold value.

Thus, counter 6, comparator 7, threshold generator 8 and random number generator 9 constitute a random period timer 5. Depending on the random number, the random period timer 5 operates as a Contend\_Fast timer or a Contend\_Slow timer.

When the random period timer 5 expires, the state machine 4 asserts Parent\_Notify state if the line state detector 3 is detecting an idle state on the cable or asserts Child\_Notify state if the line state detector is detecting a Parent\_Notify state.

During a tree identify process, a leaf node first changes its output line from idle state to Parent\_Notify state, asserting itself as a child of the corresponding branch node. If the branch node responds to this by changing its output line from idle to Child\_Notify state,

the leaf node recognizes that it is a child of the branch node and relinquishes the Parent\_Notify request. A parent-child relationship is established in this way. However, if the branch node has already asserted Parent\_Notify state after the leaf node asserted Parent\_Notify state, a collision occurs on the serial bus 2 and both signals changes to a contention-indicating state defined by the 1394 standard. Thus, the branch and leaf nodes recognize that a root contention has occurred.

Note that Parent\_Notify and Child\_Notify states are asserted in the form of continuous signals. Simultaneous presence of Parent\_Notify states on a serial bus causes the transmitted continuous signals to change their format at the point of collision and transmitted as a corrupted signal in such a way that the receiving node can recognize the occurrence of a contention. Thus, the corrupted signal is present on the transmission medium 2 until the sending node relinquishes the Parent\_Notify request. However, the format of these protocol signals is such that even though Parent\_Notify and Child\_Notify are simultaneously present on the transmission medium, the state of these signals does not change during propagation through the serial bus.

Fig. 3A shows a time sequence in which the operation of the prior art handshaking circuits of a node A (leaf) and a node B (branch) proceeds when such a contention occurs.

When the state machine 4 of node A knows that the line state detector 3 is detecting an idle state, it asserts Parent\_Notify request at time  $t = 0$ . Following a propagation delay time of 20 nanoseconds, the line state detector 3 of node B recognizes that the line state has changed to Parent\_Notify. If the state machine 4 of node B has already asserted Parent\_Notify state, node B will recognize, at time  $t = 20$  ns, that a contention has occurred, and immediately relinquishes Parent\_Notify request so that the cable changes to idle state. Thus, at time  $t = 30$  ns, node A detects the contention and relinquishes its Parent\_Notify state, changing the line to idle.

At nodes A and B, counters 6 and random number generators 9 are started at instants  $t = 30$  ns and  $t = 20$  ns, respectively. If the random numbers of nodes A and B are 0 and 1, respectively, node A is said to have a Contend\_Fast timer (260 ns) and node B a Contend\_Slow timer (600 ns). Thus, at time  $t = 290$  ns, node A asserts Parent\_Notify request again.

At time  $t = 600$  ns, the Contend\_Slow timer expires. Due to the presence of the Parent\_Notify state, node B asserts Child\_Notify state, recognizing itself as the parent of leaf node A. Node A thus recognizes itself as a child of branch node B.

However, the time difference between the Contend\_Slow and Contend\_Fast timers is determined on the condition that the maximum bus length is 4.5 meters. Therefore, if the bus length is longer than 4.5 meters, unfavorable situations might occur and root contention cannot be resolved.

For example, if the bus length is 50 meters, the propagation delay time of the bus is approximately 250 ns as shown in Fig. 3B. If the nodes A and B obtain the same random numbers as in the case of Fig. 3A, node A will recognize the occurrence of a root contention at time  $t = 430$  ns and asserts a Parent\_Notify request at time  $t = 690$  ns. Node B, on the other hand, recognizes, at time  $t = 850$  ns, that its input line is in an idle state and asserts a Parent\_Notify request. Since this change of state is recognized by node B at time  $t = 940$  ns, a root contention occurs again.

In order to overcome this problem, the present invention provides a handshaking circuit as shown in Fig. 4. This handshaking circuit comprises receive and transmit counters 10 and 11 and a comparator 12 in addition to the line state detector 3 and state machine 4 of Fig. 2.

When the state machine 4 asserts a Parent\_Notify request on the serial bus 2, it simultaneously enables the transmit counter 11 to start incrementing a transmit count (Tx). When the line state detector 3 detects a root contention, it disables the transmit counter 11 to stop the transmit count and enables the receive counter 10 to start incrementing a receive count (Rx), while informing the state machine 4 of the occurrence of a contention. State machine 4 responds to the contention by relinquishing the Parent\_Notify request. When the line state detector 3 detects an idle state, it disables the receive counter 10 to stop the receive count value, and informs the state machine 4 of this idle line state.

Thus, the counters 10 and 11 have stopped counting their clock, and supply their outputs to the comparator 12. The receive and transmit count values are compared with each other and their relative values are determined to produce a comparator output. The comparator output is supplied to the state machine 4. If the comparator output of a given node indicates that the transmit counter 11 has attained a greater value than that of receive counter 10, the state machine 4 of this node is the first to assert a Parent\_Notify request again. If the comparator output indicates otherwise (i.e., the output of transmit counter 11 is smaller than that of receive counter 10), the state machine 4 of this node may assert a Child\_Notify state in response to the output of comparator 12 or in response to a Parent\_Notify indication from the line state detector 3.

The time sequence diagram shown in Fig. 5 illustrates the operation of the handshaking circuit of Fig. 4 when a root contention arises on a serial bus with a propagation delay of 250 ns between nodes A and B which are leaf and branch nodes, respectively.

When the state machine 4 of node A recognizes that the line state detector 3 is detecting that the input line is in an idle state, it changes its output line to Parent\_Notify state at time  $t = 0$  ns, and enables its transmit counter 11 to start incrementing its count from the initial value (Tx-A = 0).

Meanwhile, at node B, the state machine 4 has

changed its output line from idle to Parent\_Notify state at time  $t = 180$  ns and has enabled its transmit counter 11 to start incrementing its count from the initial value ( $Tx-B = 0$ ).

At time  $t = 250$  ns, the line state detector 3 of node B detects a root contention, notifies the state machine 4 of this fact, enables the receive counter 10 to start incrementing the receive count from the initial value ( $Rx-B = 0$ ) and disables the transmit counter 11 so that a transmit count value  $Tx-B = 70$  is obtained. State machine 4 of node B thus relinquishes the Parent\_Notify request and the line state changes to idle.

At time  $t = 430$  ns, the line state detector 3 of node A detects the root contention, notifies the state machine 4 of this fact, enables the receive counter 10 to start incrementing its count from the initial value ( $Rx-A = 0$ ) and disables the transmit counter 11 so that a transmit count value of  $Tx-A = 430$  is obtained. State machine 4 of node A thus relinquishes the Parent\_Notify request and the line state changes to idle.

At time  $t = 500$  ns, the line state detector 3 of node A senses that its input line has changed to idle state, notifies the state machine 4 of this fact and disables the receive counter 10, producing a receive count value of  $Rx-A = 70$ . Since  $Tx-A$  is greater than  $Rx-A$  and since node A is currently detecting idle state, it is entitled to assert a Parent\_Notify request again.

At time  $t = 680$  ns, the receive counter 10 of node B is disabled, producing a receive count value of  $Rx-B = 430$ . Since,  $Tx-B$  is smaller than  $Rx-B$ , the state machine 4 of node B asserts a Child\_Notify response when it is informed, at time  $t = 750$  ns, that the state of the serial bus has changed from idle to Parent\_Notify.

As a result, a contention between nodes A and B is resolved and a parent-child relationship is established, regardless of the length of the serial bus, and node A recognizes itself as a child of branch node B and branch node B as the parent of leaf node A.

Fig. 6 shows a modified embodiment of the present invention in which parts corresponding in significance to those in Figs. 2 and 4 are marked with the same numerals as those in these figures. This modification is a combination of the handshaking circuits of the prior art and the previous embodiment and a decision circuit that determines which of these handshaking circuits is to be used depending on the transmit and receive count values relative to a predefined threshold.

For this purpose, comparators 20 and 21 are connected to the outputs of receive and transmit counters 10 and 11, respectively, for comparison with a timeout threshold value representing a delay time of 100 ns, for example, which is twice the round-trip propagation delay (50 ns) of a bus length 4.5 meters. When one of the receive and transmit counters 10 and 11 exceeds the timeout threshold value, the corresponding one of comparators 20 and 21 produces a binary 1 at the corresponding input of an OR gate 22. Thus, when both counters do not exceed the threshold value both com-

parators 20 and 21 produce a binary 0 and hence the output of OR gate 22 is at binary 0. The output of OR gate 22 is connected to a selector 23, to which the output comparators 7 and 12 are also connected.

When the output of OR gate 22 is a binary 0, the selector 23 selects and couples the output of comparator 7 to the state machine 4. When the output of OR gate 22 is a binary 1, the selector 23 selects and couples the output of comparator 12 to the state machine 4. State machine 4 is supplied with the output of OR gate 22 via a line 25 to indicate in which mode the state machine is to be operated.

When either of the receive and transmit counters do not exceeds the threshold of 100 ns (thus, the output of OR gate is binary 0) and the comparator 7 produces a timeout signal, the state machine 4 is signaled by the output of OR gate 22 via line 25 that it should respond to the output of comparator 7 for operating in the random timer mode of Fig. 2. When one of the receive and transmit counters exceeds the threshold of 100 ns (thus, the output of OR gate is binary 1) and the comparator 12 produces an output indicating which of the receive and transmit counters is greater, the state machine 4 is signaled by the output of OR gate 22 via line 25 that it should respond to the output of comparator 12 for operating in the counter mode of Fig. 4.

In summary, the elements 20, 21, 22 and 23 constitute a decision circuit 24 that supplies the state machine 4 with a first decision output to allow it to operate in the random time mode when the transmit and receive count values are both lower than the 100-ns threshold and with a second decision output to allow it to operate in the counter mode when one of the transmit and receive count values is greater than the threshold. State machine 4 responds to the first decision output for asserting the Parent\_Notify state on the serial bus if the random period timer 5 expires during the presence of idle on the serial bus and asserting the Child\_Notify state if the random period timer 5 expires during the presence of the Parent\_Notify state. State machine 4 responds to the second decision output for asserting the Parent\_Notify state when the transmit count value is greater than the receive count value and asserting the Child\_Notify state when the transmit count value is smaller than the receive count value.

As illustrated in Fig. 7, if node B asserts Parent\_Notify state 180 ns after node A initially asserted Parent\_Notify in a manner similar to Fig. 5, a contention occurs. If the serial bus 2 is 50 meters long (corresponding to a single-trip propagation delay of 250 ns), the receive counter 10 of node B reaches the threshold value of the comparator 20 at time  $t = 350$  ns, which, as described above, corresponds to the point of determination for node B between the random time mode and the counter mode, as illustrated in Fig. 7.

Since the transmit counter 11 of node A has already exceeded the threshold value of comparator 21 when its line state detector 3 detects the contention at time  $t =$

430 ns, and the transmit counter 11 of node B has already exceeded the threshold value of comparator 21 when its line state detector 3 detects the idle state at time  $t = 680$  ns, nodes A and B operate in the counter mode to resolve the contention in a manner similar to Fig. 5. It is seen that if the round-trip propagation delay of the serial bus is 40 ns, for example, the handshaking circuit operates in the random time mode to resolve the contention as shown in Fig. 3A.

Fig. 8 shows another embodiment of the handshaking circuit in which the prior art threshold generator 8 of Fig. 2 is modified so that it comprises a register 40 for holding a fixed lower threshold value corresponding to the Contend\_Fast time of 260 ns, a register 41 for holding a fixed higher threshold value corresponding to the Contend\_Slow time of 600 ns, and a register 42 for storing a variable threshold value as an additional Contend\_Slow timer value. An adder 43 is provided for summing the contents of registers 41 and 42. One of the outputs of adder 43 and register 40 is selected by a selector 44 according to the output of random number generator 9. The output of selector 44 is supplied as a random threshold value to the comparator 7. The fixed value of register 41 corresponds to the nominal length of the serial bus. The value of register 42 is manually set so that it corresponds to the distance by which the length of the serial bus exceeds the nominal bus length.

If the additional threshold value of register 42 corresponds to a delay time of 180 ns and the serial bus has a length of 50 meters (where the round-trip propagation delay is approximately 500 ns), the line state detector 3 of node B detects Parent\_Notify state within the extended time window of 180 ns, as shown in Fig. 9, provided that the same random numbers are adopted by nodes A and B as those in Fig. 3B. As a result, node B asserts Child\_Notify state when its line state detector 3 senses a Parent\_Notify state at time  $t = 1030$  ns. Root contention on a serial bus longer than the nominal value can be resolved by adjusting the value of register 42.

Fig. 10 is a block diagram illustrating another embodiment of this invention in which repeaters 50 and 51 (Fig. 11) are provided at the opposite end of the serial bus 2 and used in combination of handshaking circuits (Fig. 12) of nodes 52 and 53.

As shown in detail in Fig. 11, each repeater includes a format converter 60 interposed between the serial bus 2 with which it is connected to the handshaking circuit of the local node and a cable 54 with it is connected to a remote repeater. Converter 60 performs conversion between the 1394 standard format on the serial bus 2 and the usual bipolar format on the cable 54.

A Root\_Contention detector 61 is provided for detecting a root contention on the serial bus 2 to produce a contention presence signal. Idle detectors 62 and 64 are connected to the serial bus 2 and the high-speed cable 54, respectively, to detect an idle state and produce an output when they are enabled by the con-

tention presence signal. A root contention notify circuit 63 is provided to respond to the output of idle detector 62 for setting the serial bus 2 at Root\_Contention\_Notify state for the purpose of locking the handshaking circuit of the local node in an inactive state. At the same time, the notify circuit 63 supplies a disable signal via an OR gate 71 to the converter 60 to prevent it from passing the signal from the remote repeater to the local node when the RC notify circuit 63 is supplying a signal to the local node, instead of the signal from the remote repeater. The output of root contention detector 61 is also applied as a disable signal via the OR gate 71 to the converter 60 to prevent it from passing the remote signal to the local node when a root contention is detected.

A Parent\_Notify detector 65 is connected to the input line of cable 54 to disable the RC notify circuit 63 when the line state changes to Parent\_Notify after the detector 65 is enabled by the output of idle detector 64.

Counter 66 starts incrementing a clock count in response to the contention presence signal from detector 61. Random number generator 69 is also responsive to the contention presence signal to produce a random number of either 1 or 0. Threshold generator 68 produces a lower threshold value corresponding to 260 ns (Contend\_Fast). According to this embodiment, this threshold generator has a higher threshold value corresponding to 780 ns (Contend\_Slow) which is larger than the higher random threshold value specified by the 1394 standard. Depending on the output of random number generator 69, the threshold generator 68 selects one of the threshold values. Comparator 67 produces an output indicating the value of the clock count relative to the threshold value and disables the state machine 63 when the clock count exceeds the threshold value. Elements 66 to 69 thus constitute a random period timer 70.

Fig. 12 shows the handshaking circuit used by the nodes 52 and 53. As illustrated, this handshaking circuit is similar to that shown in Fig. 2 except that the line state detector 3 responds to the Root\_Contention\_Notify state on the serial bus 2 from the local repeater for locking the state machine 4 so that it is prevented from responding to a state indication signal from the line state detector 3 at the time the random period timer 5 expires.

The following is a description of the root contention resolution procedure of the repeaters and handshaking circuits of Figs. 11 and 12 with the aid of the time sequence diagram of Fig. 13 by assuming that node A and repeater A determine a period of 260 ns, repeater B determines a period of 780 ns and node B determines a period of 600 ns.

At time  $t = 0$ , node A (leaf node) asserts Parent\_Notify state, which is converted by repeaters A and B and detected by node B at time  $T_2$  as a root contention since node B has already asserted Parent\_Notify at time  $T_0$ . Thus, repeater B detects a root contention at time  $T_1$  and starts its counter 66 and

node B changes the local serial bus 2 to an idle state at time T2 and enables its counter 6. Repeater B changes the local serial bus 2 to Root\_Contention\_Notify state (instead of the corrupted Parent\_Notify state) at time T3 when its idle detector 62 senses the idle state on the serial bus 2.

At time T4, the Parent\_Notify state asserted at T0 is sensed by repeater A as a root contention, and relayed to node A, which also senses it at time T5 as a root contention and relinquishes the Parent\_Notify state and sets the serial bus to an idle state, while starting its counter 6. Repeater A detects this idle state at time T6 and changes the input line of node A to Root\_Contention\_Notify state (instead of the corrupted Parent\_Notify state).

At time T7, the counter 66 of repeater A reaches the random period of 260 ns, causing the contention notify circuit 63 to relinquish the R\_C\_Notify state, thus setting the serial bus to idle. At time T8, the counter 6 of node A reaches the Contend\_Fast period of 260 ns.

Due to the occurrence of the idle state at its input line at time T8, node A, which has been locked from time T6, is now unlocked and asserts a Parent\_Notify state again as a child of node B.

At time T9, the counter 6 of node B reaches the Contend\_Slow period of 600 ns. Since node B has been in a locked state from time T3' due to the presence of the R\_C\_Notify state, its state machine 4 does not change its Parent\_Notify state although its Contend\_Slow period already expired at time T9.

At time T10, the Parent\_Notify detector 65 of repeater B detects the Parent\_Notify state and disables the contention notify circuit 63 so that the serial bus 2 of node B changes from R\_C\_Notify to Parent\_Notify state, which the node B senses at time T11. Thus, node B responds to this Parent\_Notify state and asserts a Child\_Notify state as a parent of node A. At time T12, the Contend\_Slow period of repeater B expires when its counter 66 reaches the timeout threshold of 780 ns.

## Claims

1. A handshaking circuit for establishing a particular relationship between nodes interconnected by a transmission medium, comprising:

a state detector (3) for monitoring said transmission medium (2) to detect a predetermined first state, a predetermined second state, a contention when predetermined first states are simultaneously asserted on said transmission medium, and an idle state when either of said first and second states is not asserted on the transmission medium;  
a state machine (4) connected to the state detector (3) for asserting the first state on said transmission medium to initiate a handshaking process and relinquishing the first state when

the contention is detected by the state detector;

a receive counter (10) connected to the state detector (3) for continuously incrementing a receive count value from the instant the contention is detected and the instant said idle state is detected;

a transmit counter (11) connected to the state detector (3) and the state machine (4) for continuously incrementing a transmit count value from the instant the first state is detected to the instant the contention is detected; and  
a comparator (12) for comparing the receive count value with the transmit count value, said state machine (4) asserting the first state on the transmission medium again when the transmit count value is greater than the receive count value and asserting said second state on the transmission medium when the transmit count value is smaller than the receive count value.

2. A handshaking circuit as claimed in claim 1, wherein said state machine (4) is arranged to assert the second state on the transmission medium at the instant the first state is detected again by the state detector.
3. A handshaking circuit for establishing a particular relationship between nodes interconnected by a transmission medium, comprising:

a state detector (3) for monitoring said transmission medium (2) to detect a predetermined first state, a predetermined second state, a contention when predetermined first states are simultaneously asserted on said transmission medium, and an idle state when either of said first and second states is not asserted on said transmission medium;

a state machine (4) connected to the state detector for asserting the first state on the transmission medium to initiate a handshaking process and relinquishing the first state when the contention is detected by the state detector;

a receive counter (10) connected to the state detector (3) for continuously incrementing a receive count value from the instant the contention is detected and the instant said idle state is detected;

a transmit counter (11) connected to the state detector (3) and the state machine (4) for continuously incrementing a transmit count value from the instant the first state is detected to the instant said contention is detected;

a comparator (12) for comparing the receive count value with the transmit count value,



decision means (24) for producing a first decision output when the transmit and receive count values are lower than a predefined threshold and a second decision output when one of said transmit and receive count values is greater than the predefined threshold; and  
 a random period timer (5) arranged to be activated in response to the detection of the contention by the state detector (3);

wherein said state machine (4) is arranged to:

respond to the first decision output for asserting the first state on the transmission medium if the random period timer (5) expires during the absence of the first and second states on the transmission medium and asserting the second state on the transmission medium if the random period timer (5) expires during the presence of the first state on the transmission medium; and  
 respond to the second decision output for asserting the first state on the transmission medium when the transmit count value is greater than the receive count value and asserting the second state on the transmission medium when the transmit count value is smaller than the receive count value.

4. A handshaking circuit for establishing a particular relationship between nodes interconnected by a transmission medium, comprising:

a random number generator (9) for producing a random number;  
 a threshold generator (8) having a first threshold value, a second threshold value, an adjustable value, means (43) for summing the second threshold value and the adjustable value to produce a summed threshold value greater than the first threshold value, and means (44) for selecting of the first threshold value and the summed threshold value depending on said random number;  
 a state detector (3) for detecting predetermined first and second states and a contention between said nodes on the transmission medium (2);  
 a counter (6) for continuously incrementing a clock count in response to the detection of the contention by the state detector (3);  
 a comparator (7) for comparing the clock count of the counter (6) with the selected threshold value and producing an output when the clock count equals the selected threshold value; and  
 a state machine (4) connected to the state detector (3) for asserting the first state on the

transmission medium to initiate a handshaking process and relinquishing the first state when the contention is detected, and asserting the first state on the transmission medium again if the comparator (7) produces said output during the absence of the first and second states on the transmission medium and asserting the second state on the transmission medium if the comparator (7) produces said output during the presence of the first state on the transmission medium.

5. A communication system comprising a handshaking circuit connected through a repeater to a remote site via a transmission medium:

the handshaking circuit comprising:

a state detector (3) for monitoring said transmission medium (2) to detect a predetermined first state, a predetermined second state, an idle state and a contention when predetermined first states are simultaneously asserted on said transmission medium; and  
 a random period timer (5) responsive to the detection of the contention by the state detector for continuously incrementing a clock count value;  
 a state machine (4) connected to the state detector (3) for asserting the first state on the transmission medium to initiate a handshaking process and relinquishing the first state when the contention is detected, asserting the first state on the transmission medium again if the random period timer (5) expires during the absence of the first or second state on the transmission medium and asserting said second state on the transmission medium if the random period timer (5) expires during the presence of the first state on the transmission medium,

the repeater comprising:

a two-way format converter (60) connected in said transmission medium so that the transmission medium is divided into a local side (2) and a remote side (54);  
 a contention detector (61) for detecting a contention on the local side of said transmission medium;  
 a first idle detector (62) responsive to the detection of the contention for detecting the idle state on the local side of the transmission medium;  
 notification means (63) responsive to the detection of the idle state by the first idle detector (62) for asserting a third state on the local side of the transmission medium and disabling



said converter (60);

a second idle detector (64) responsive to the detection of said contention for detecting the idle state on the remote side of the transmission medium;

a state detector (65) responsive to the detection of the idle state by the second idle detector (64) for monitoring the remote side of the transmission medium for disabling the notification means (63) when the first state is detected on the remote side of the transmission medium; and

a random period timer (70) for disabling the notification means (63) when a random interval from the time of detection of said contention expires,

the state detector (3) of the handshaking circuit being arranged to detect the third state on the local side of the transmission medium, and

the state machine (4) of the handshaking circuit being arranged to be disabled in response to the detection of the third state so that the state machine (4) is prevented from responding to the expiration of said random period timer (5) of the handshaking circuit as long as the third state is detected on the local side of the transmission medium.

6. A method for resolving a contention that occurs on a transmission medium between two network nodes, using transmit and receive counters, the method comprising:

a) monitoring said transmission medium (2) to detect a predetermined first state, a predetermined second state, an idle state when either of said predetermined first and states is not present on the transmission medium, and a contention between said network nodes when predetermined first states are simultaneously present on the transmission medium;

b) asserting the first state on said transmission medium to initiate a handshaking process and enabling the transmit counter (11) to start incrementing a transmit count value;

c) when said contention is detected, relinquishing the first state, disabling the transmit counter (11) to stop incrementing the transmit count value and enabling the receive counter (10) to start incrementing a receive count value;

d) disabling the receive counter (10) to stop incrementing the receive count value when the idle state is detected;

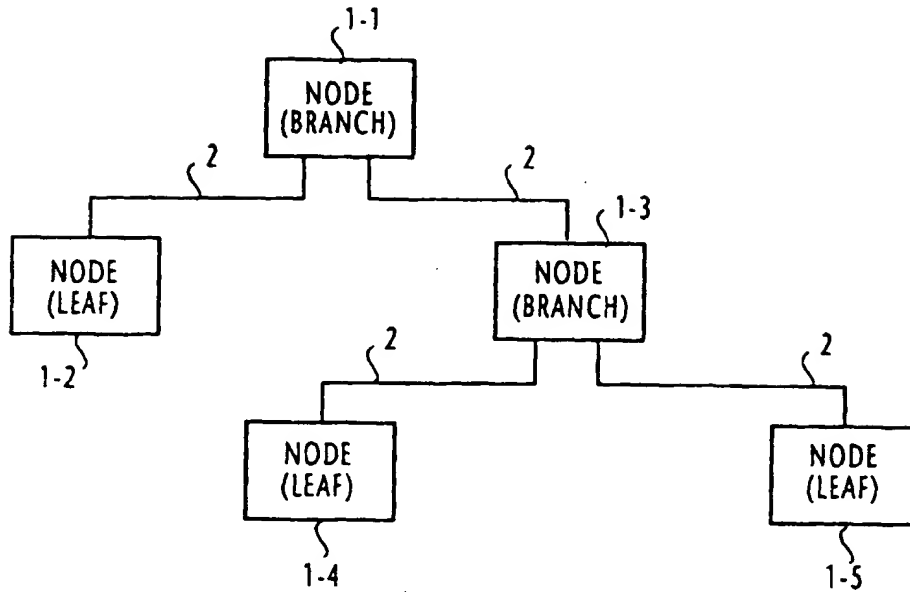
e) comparing the receive count value with the transmit count value; and

f) asserting the first state on the transmission medium again when the transmit count value is greater than the receive count value and

asserting the second state on the transmission medium when the transmit count value is smaller than the receive count value.

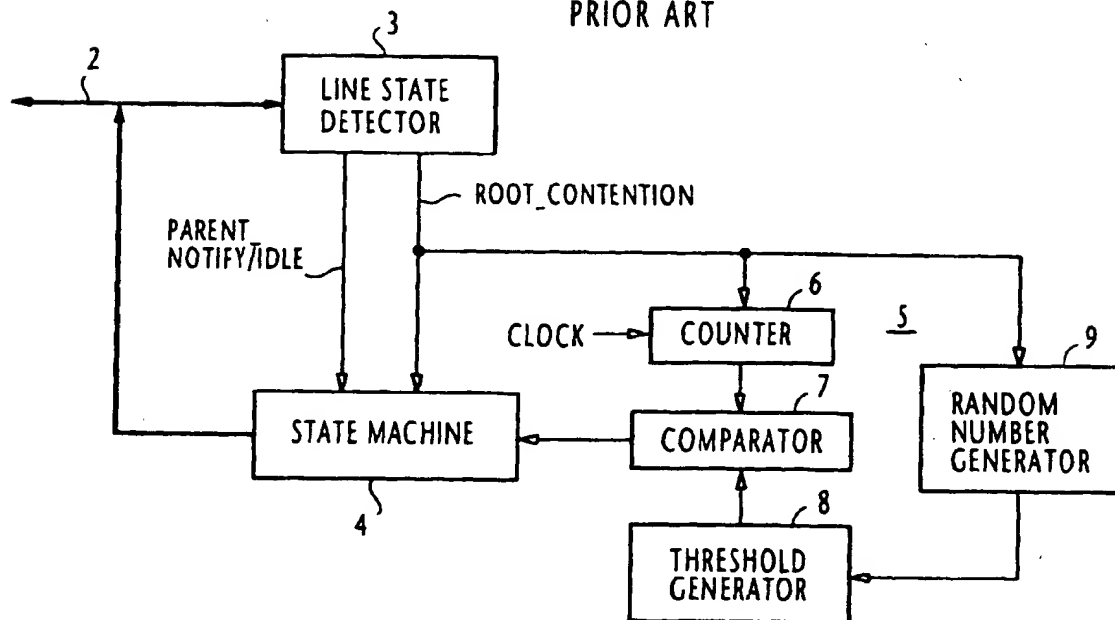
7. The method of claim 6, wherein the step (g) comprises asserting said second state on the transmission medium at the instant the first state is detected again on the transmission medium.

**FIG. 1**

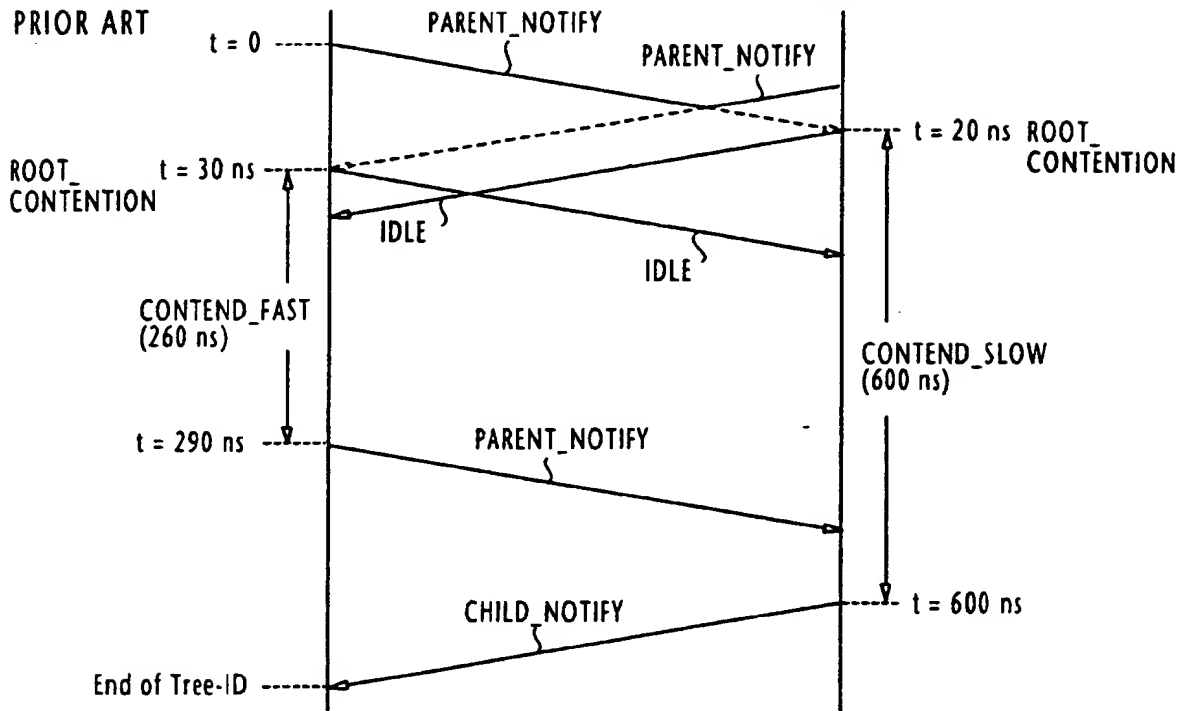


**FIG. 2**

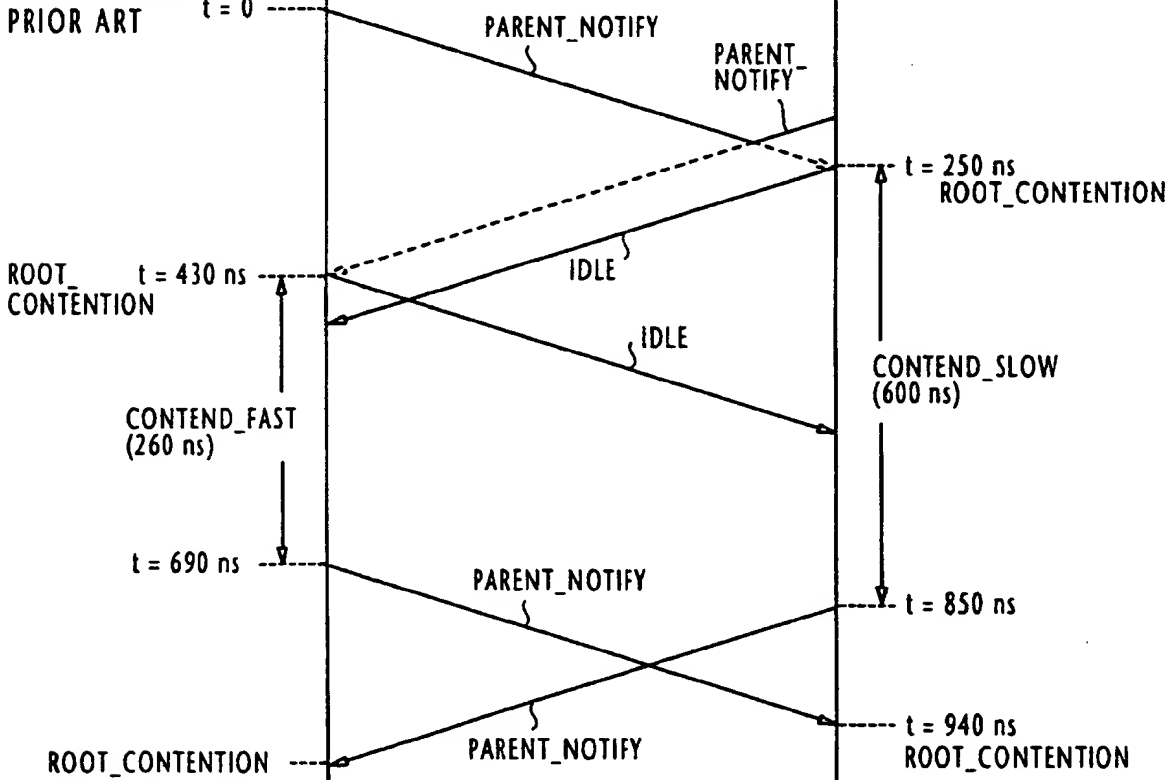
PRIOR ART



**FIG. 3A**



**FIG. 3B**



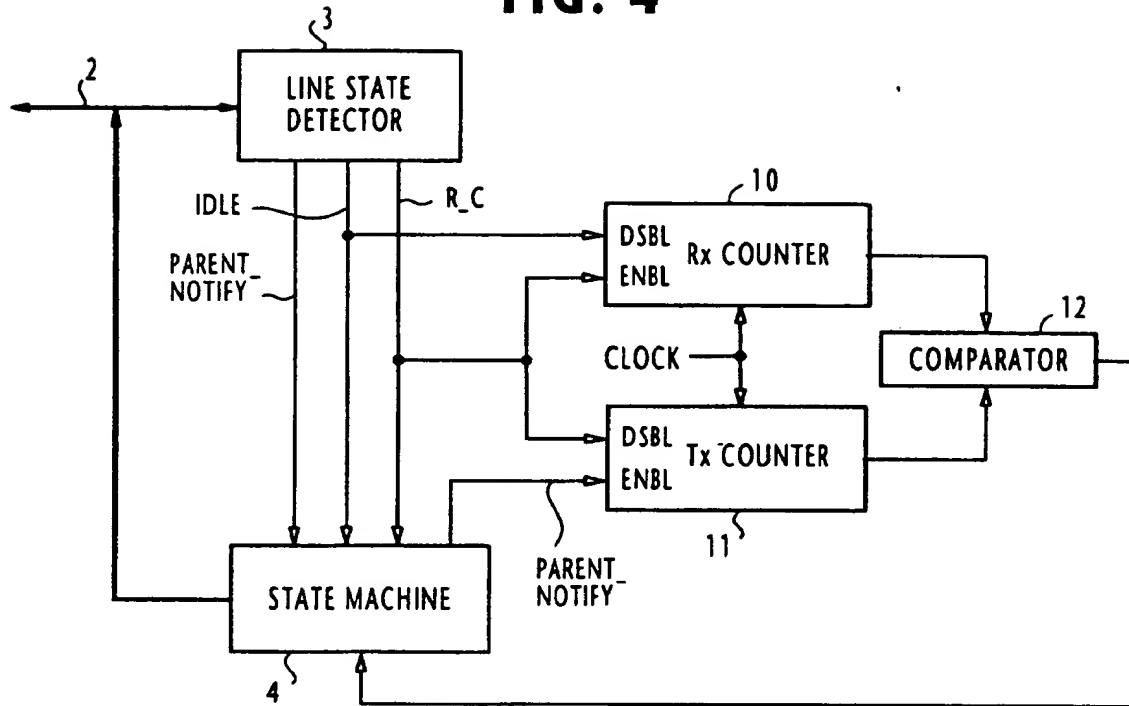
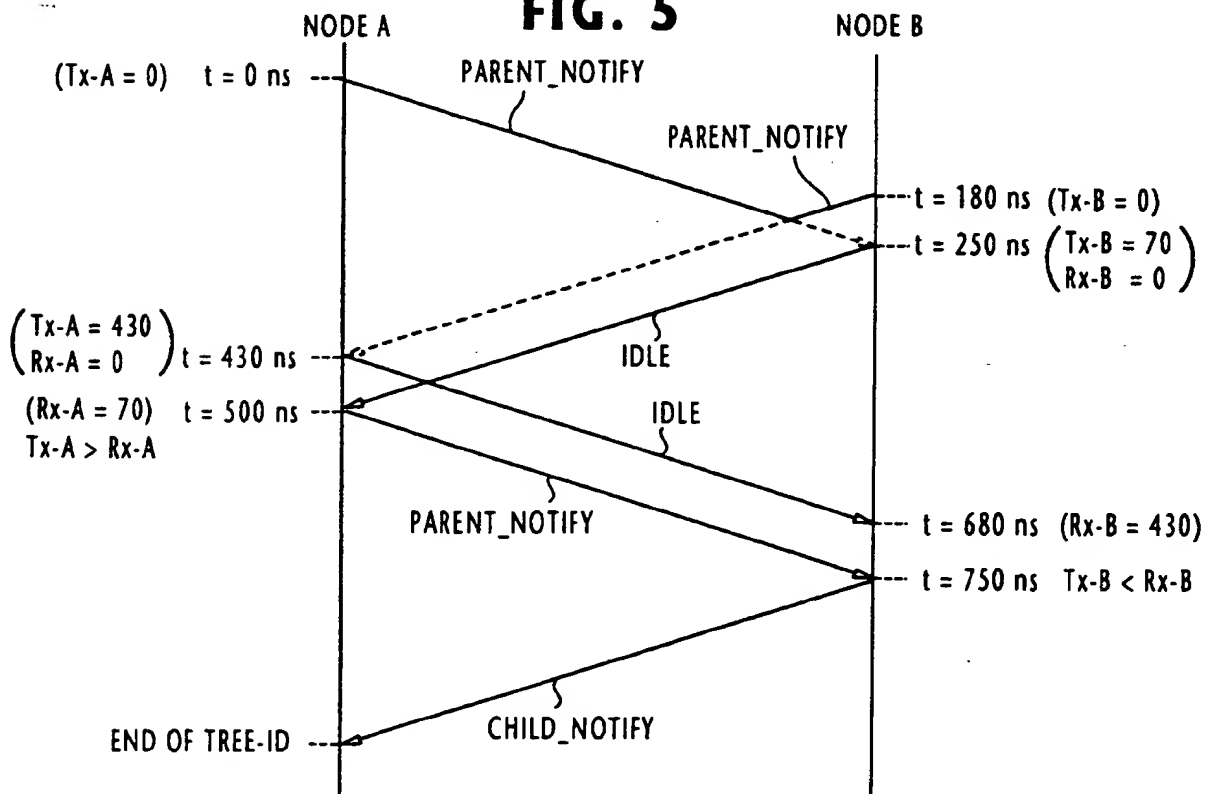
**FIG. 4****FIG. 5**

FIG. 6

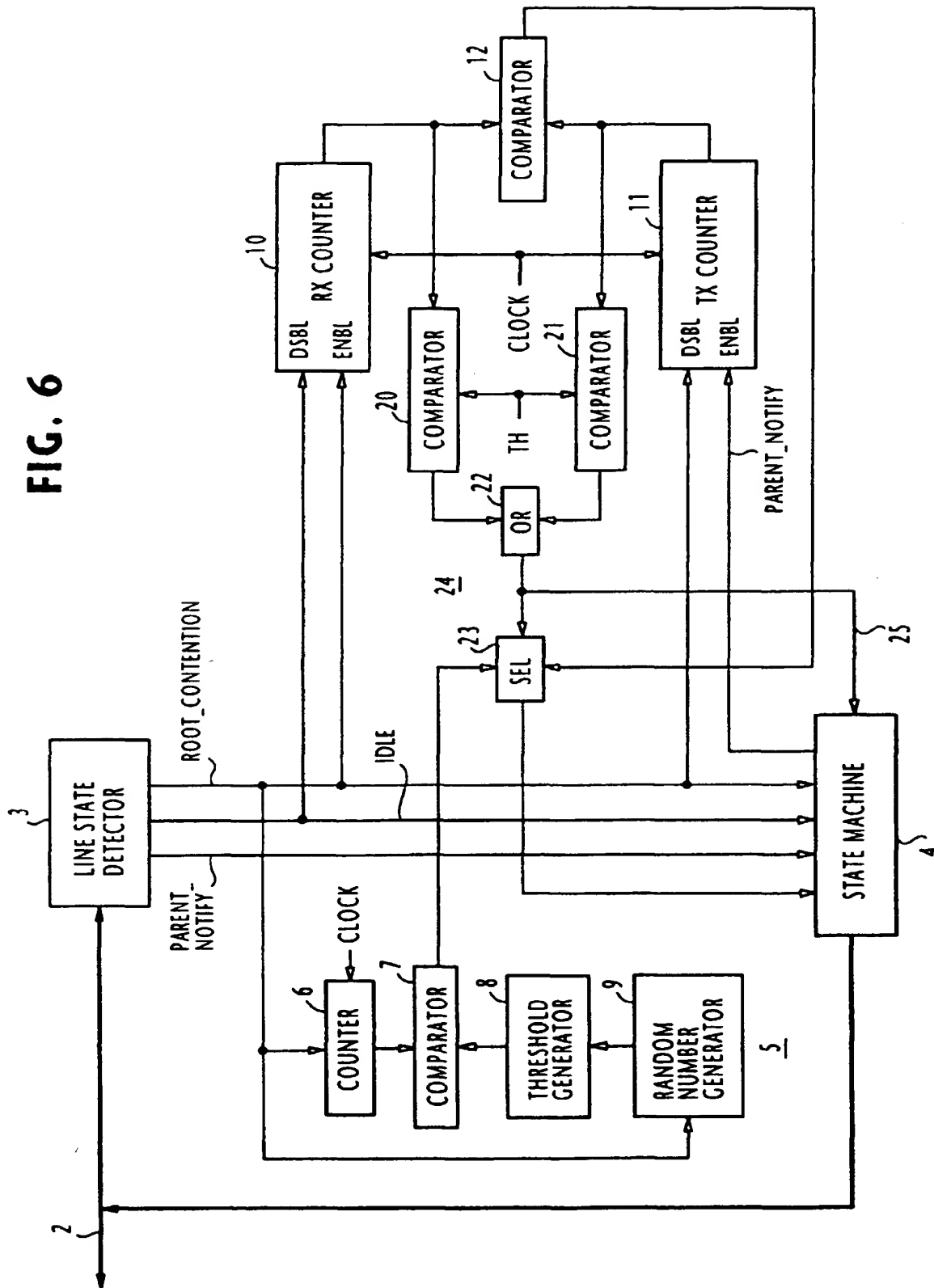


FIG. 7

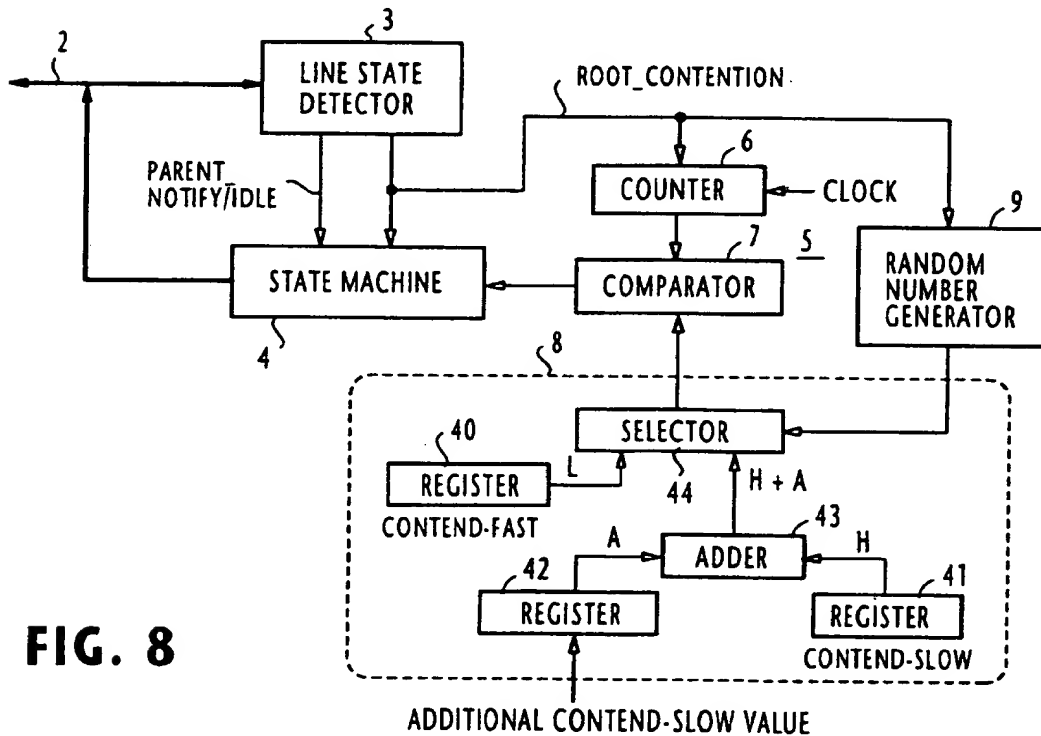
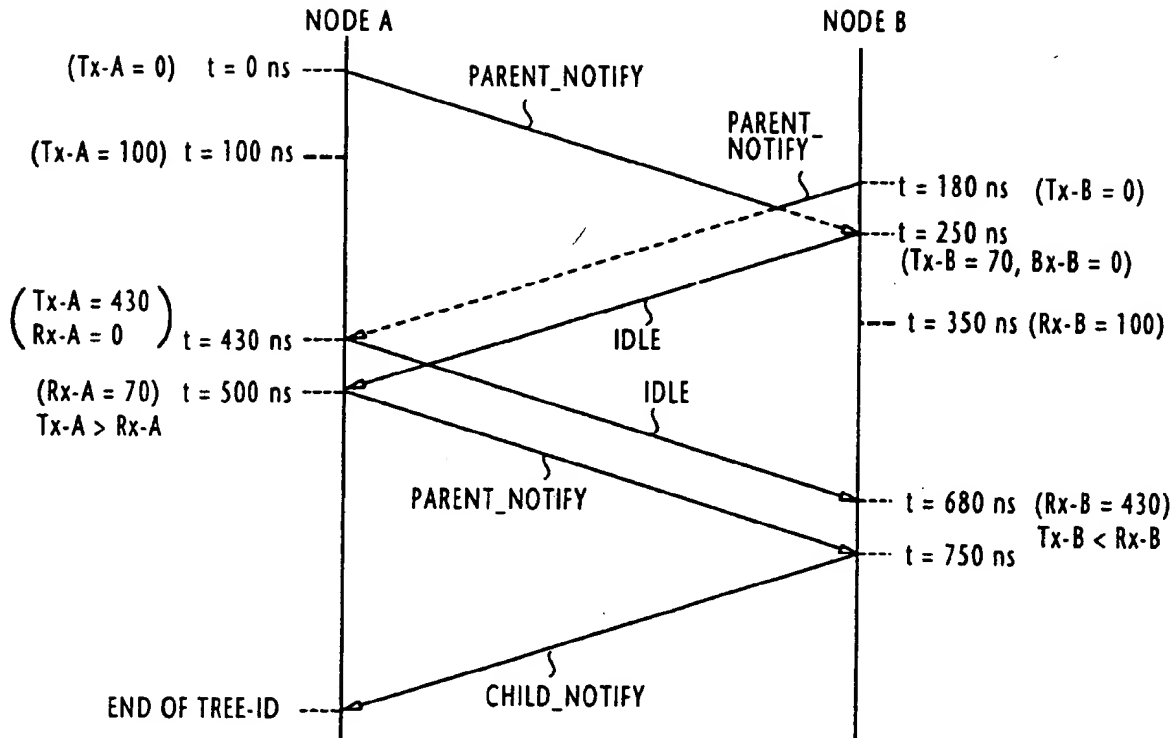
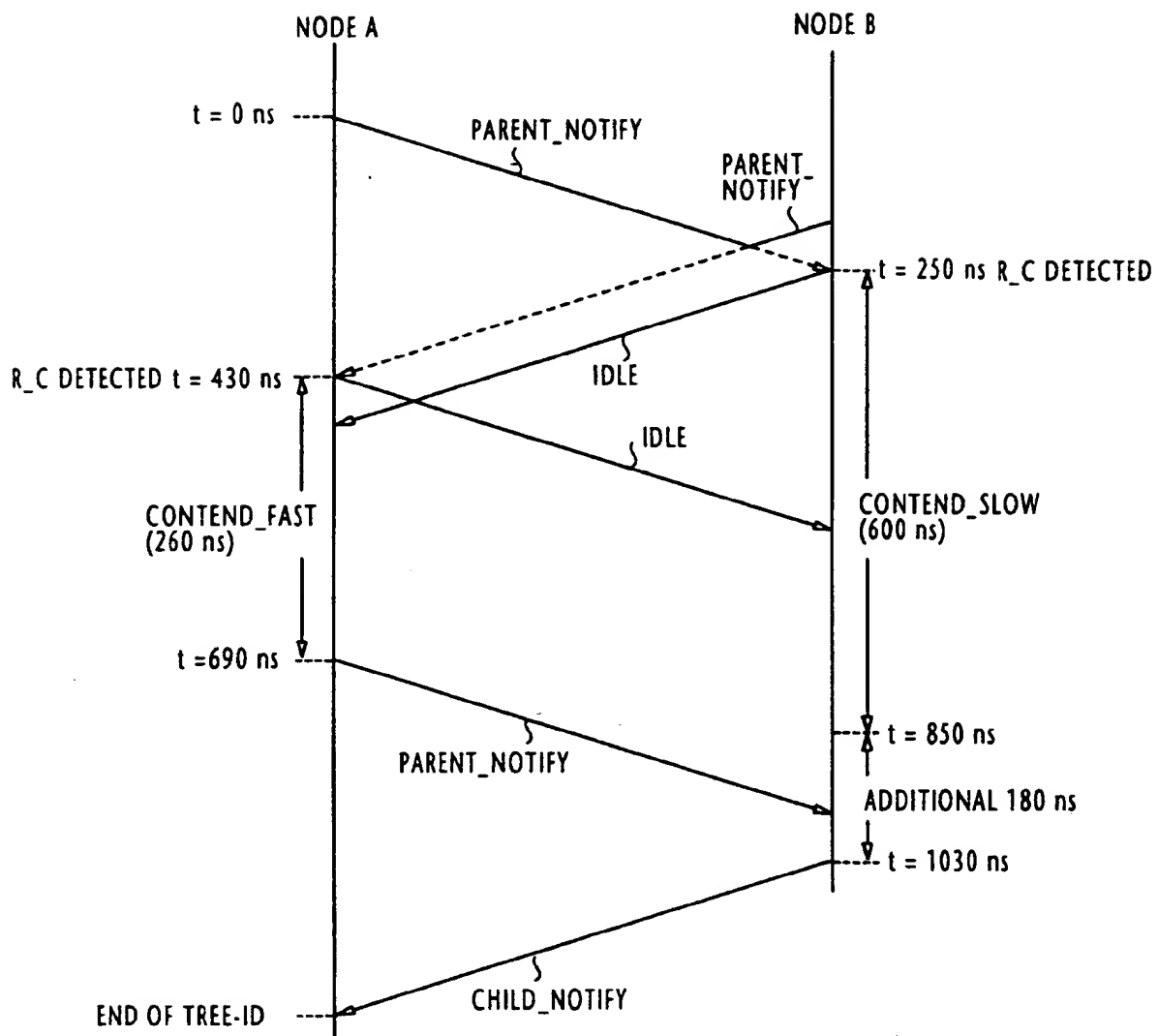
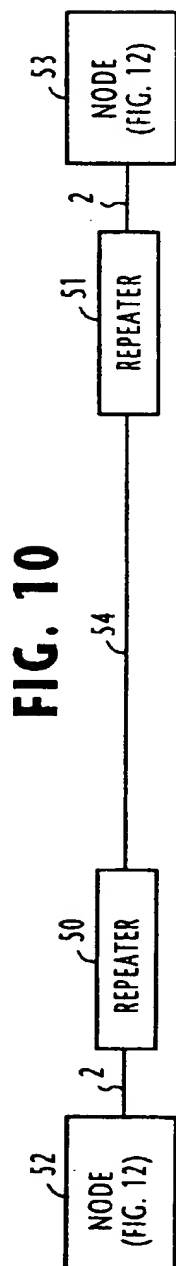


FIG. 8

**FIG. 9**







**FIG. 11**

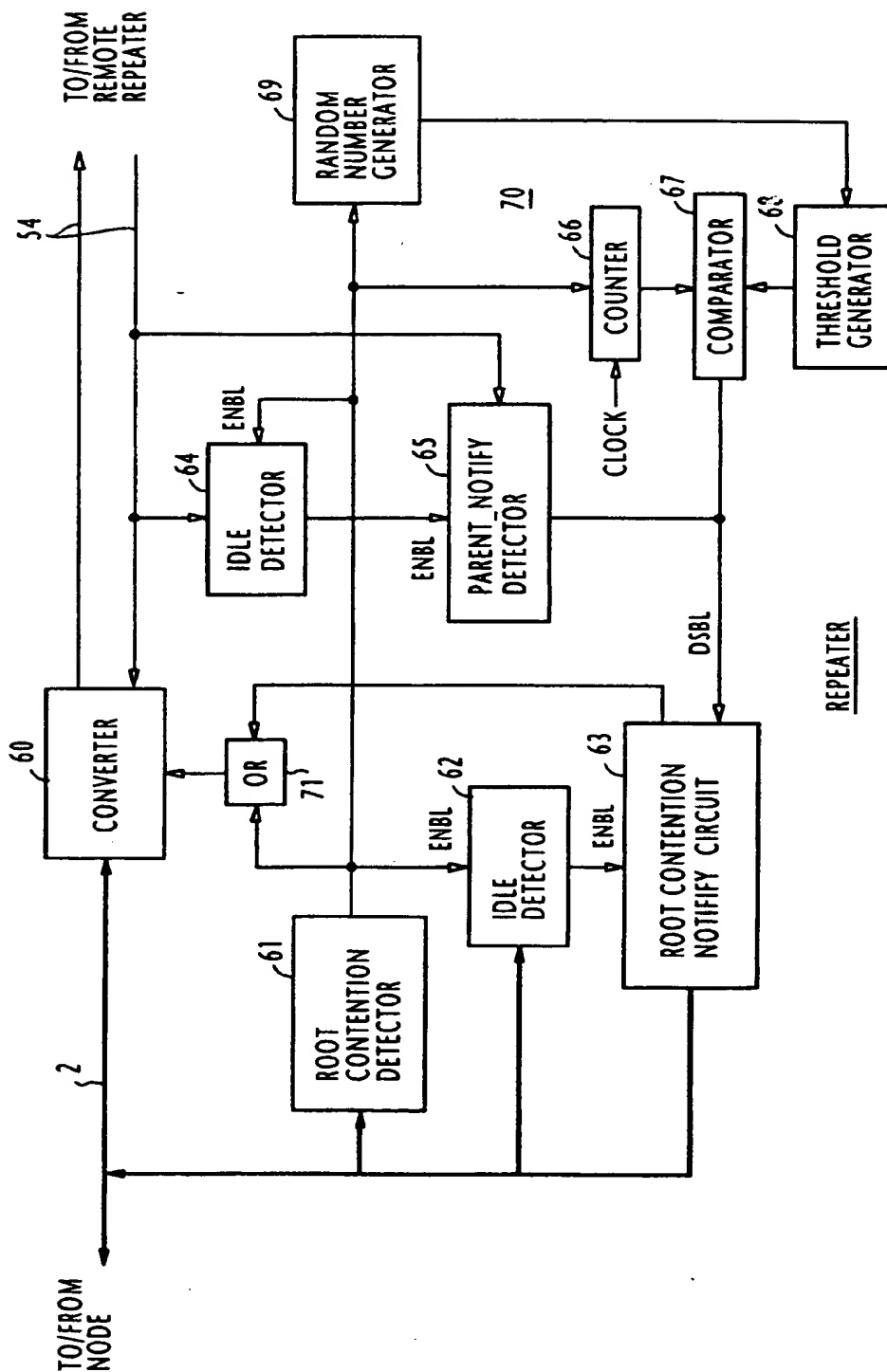


FIG. 12

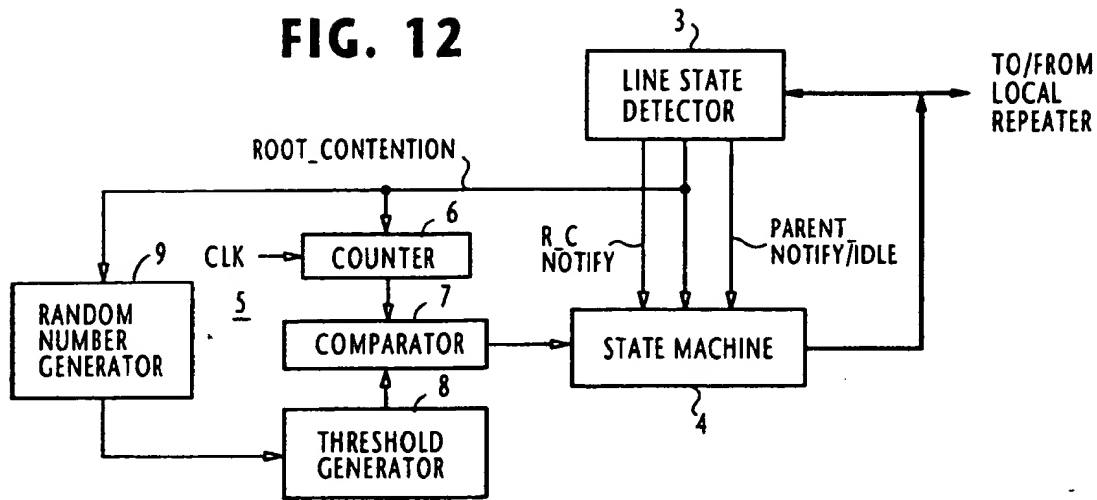
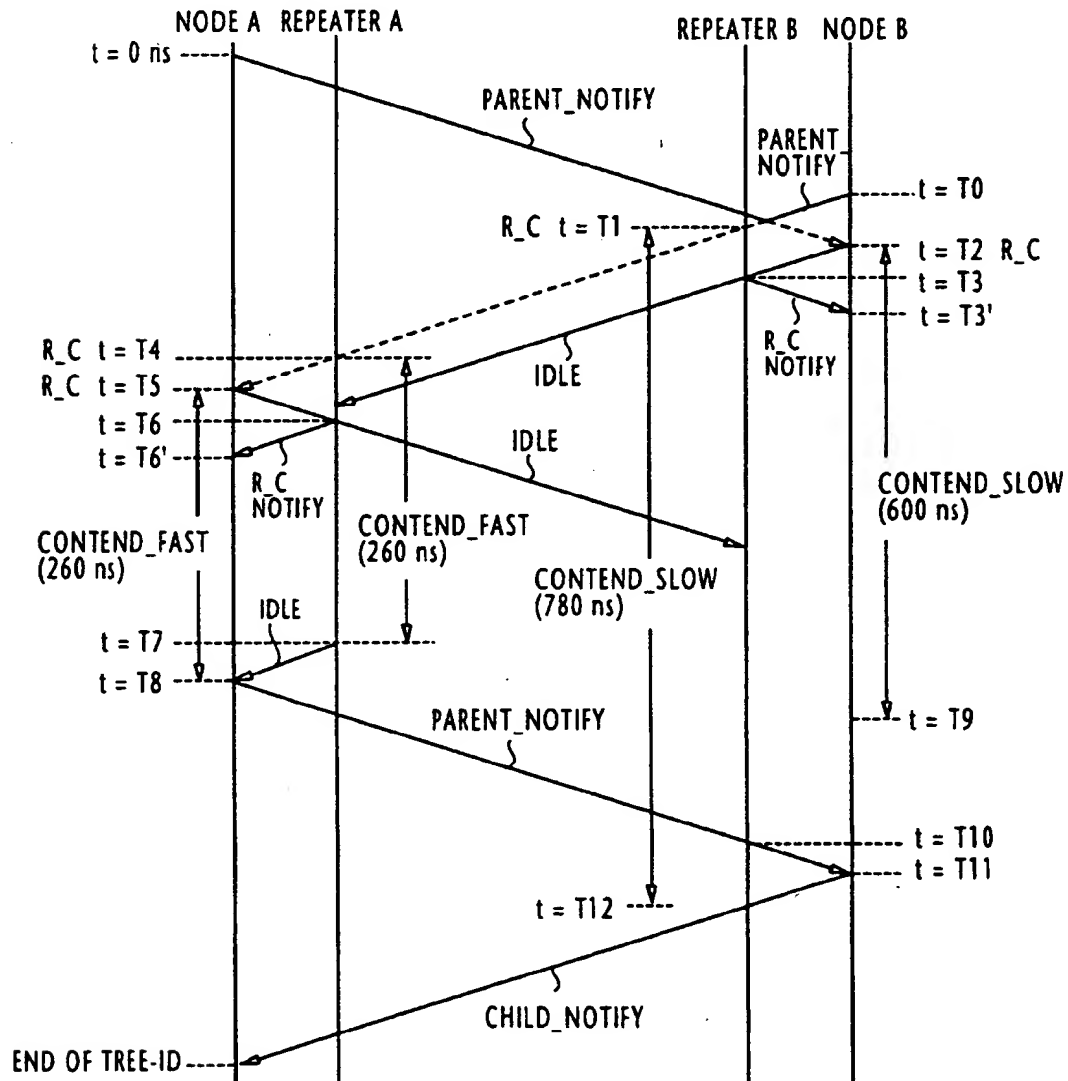
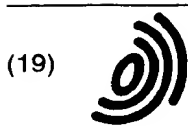


FIG. 13





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**Minato-ku, Tokyo (JP)**

(30) Priority: **26.03.1997 JP 7318797**

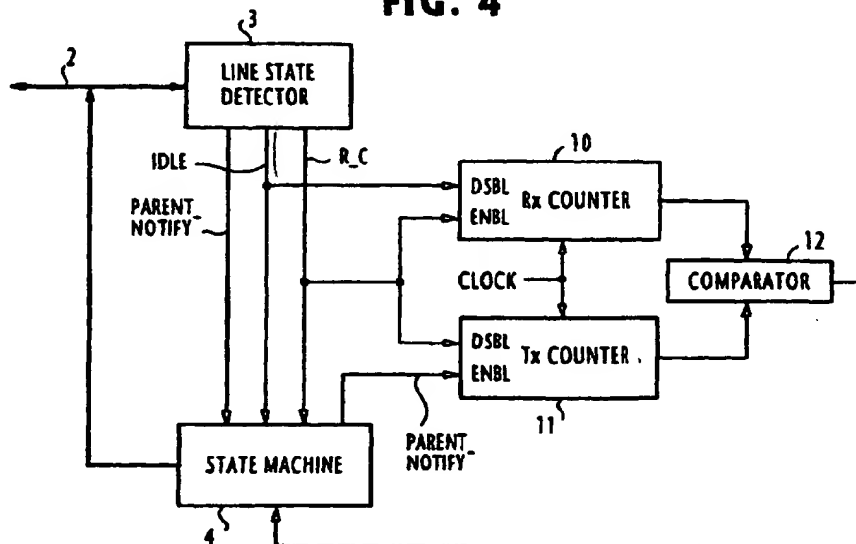
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(54) **Handshaking circuit for resolving contention on a transmission medium regardless of its length**

(57) For establishing a parent-child relationship between adjacent network nodes interconnected by a transmission medium, a handshaking circuit comprises a state detector for monitoring the transmission medium to detect predetermined first and second states and a contention that occurs when predetermined first states are simultaneously present. A state machine asserts the first state for initiating a handshaking process and relinquishes this state when a contention is detected. A re-

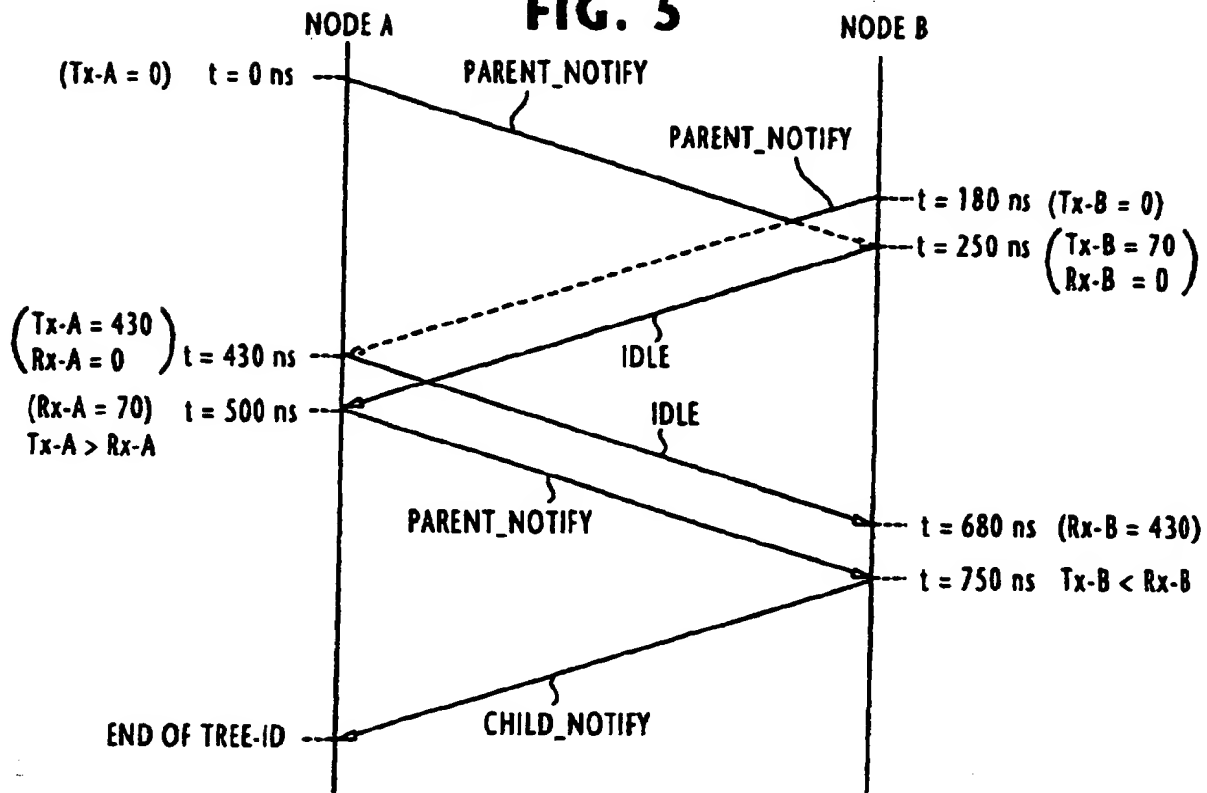
ceive count value is continuously incremented from the instant the contention is detected to the instant the transmission medium changes to an idle state. A transmit count value is continuously incremented from the instant the first state is detected to the instant the contention is detected. When the transmit count value is greater than the receive count value, the state machine asserts the first state again and when the transmit count value is smaller than the receive count value, it asserts the second state.

**FIG. 4**



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**FIG. 5**





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# EUROPEAN SEARCH REPORT

Application Number  
EP 98 10 5398

DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.6)
A	US 5 293 375 A (MOORWOOD CHARLES A ET AL) 8 March 1994 (1994-03-08) * abstract * * column 7, line 31 - column 12, line 4 * * column 17, line 6 - column 18, line 39 *	5	
A	MERAKOS L ET AL: "Interconnection of CSMA/CD LANs via an N-port bridge" INFOCOM '89. PROCEEDINGS OF THE EIGHTH ANNUAL JOINT CONFERENCE OF THE IEEE COMPUTER AND COMMUNICATIONS SOCIETIES. TECHNOLOGY: EMERGING OR CONVERGING, IEEE OTTAWA, ONT., CANADA 23-27 APRIL 1989, WASHINGTON, DC, USA, IEEE COMPUT. SOC. PR, US, 23 April 1989 (1989-04-23), pages 28-37, XP010015339 ISBN: 0-8186-1920-1 * abstract * * par. 2*	5	
The present search report has been drawn up for all claims			TECHNICAL FIELDS SEARCHED (Int.Cl.6)
Place of search <b>MUNICH</b>		Date of completion of the search <b>26 March 2003</b>	Examiner <b>Falò, L</b>
<p><b>CATEGORY OF CITED DOCUMENTS</b></p> <p>X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document</p> <p>T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons &amp; : member of the same patent family, corresponding document</p>			

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## EUROPEAN SEARCH REPORT

Application Number  
EP 98 10 5398

DOCUMENTS CONSIDERED TO BE RELEVANT			
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Y	US 4 547 850 A (GENMA HIDEAKI) 15 October 1985 (1985-10-15) * abstract *	4	
A	* column 1, line 1 - line 44 * * column 2, line 28 - line 39 * * column 5, line 11 - column 6, line 29 *	1-3,6,7	
Y	IEEE STANDARD FOR A HIGH PERFORMANCE SERIAL BUS (IEEE STD. 1394-1995), 'Online! 1996, XP002236156 ISBN: 1-55937-583-3 Retrieved from the Internet: <URL:http://1eeexplore.1eee.org/xp1/confer ences.jsp> 'retrieved on 2003-03-26!	4	
A	* page 102 *	1-7	TECHNICAL FIELDS SEARCHED (Int.Cl.6)
A	US 4 707 829 A (PENDSE SUDHIR B) 17 November 1987 (1987-11-17) * abstract * * column 2, line 59 - column 3, line 26 *	1-3,6,7	H04L
A	US 4 598 285 A (HOSHEN JOSEPH) 1 July 1986 (1986-07-01) * abstract * * claim 1 *	1-3,6,7	
A	US 4 430 651 A (BRYANT DAVID M ET AL) 7 February 1984 (1984-02-07) * column 5, line 26 - column 6, line 59 *	4	
The present search report has been drawn up for all claims			
Place of search <b>MUNICH</b>		Date of completion of the search <b>26 March 2003</b>	Examiner <b>Falò, L</b>
<p><b>CATEGORY OF CITED DOCUMENTS</b></p> <p>X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : intermediate document</p> <p>T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons &amp; : member of the same patent family, corresponding document</p>			

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Application Number  
EP 98 10 5398

### CLAIMS INCURRING FEES

The present European patent application comprised at the time of filing more than ten claims.

- ☐ Only part of the claims have been paid within the prescribed time limit. The present European search report has been drawn up for the first ten claims and for those claims for which claims fees have been paid, namely claim(s):
- ☐ No claims fees have been paid within the prescribed time limit. The present European search report has been drawn up for the first ten claims.

### LACK OF UNITY OF INVENTION

The Search Division considers that the present European patent application does not comply with the requirements of unity of invention and relates to several inventions or groups of inventions, namely:

see sheet B

- ☒ All further search fees have been paid within the fixed time limit. The present European search report has been drawn up for all claims.
- ☐ As all searchable claims could be searched without effort justifying an additional fee, the Search Division did not invite payment of any additional fee.
- ☐ Only part of the further search fees have been paid within the fixed time limit. The present European search report has been drawn up for those parts of the European patent application which relate to the inventions in respect of which search fees have been paid, namely claims:
- ☐ None of the further search fees have been paid within the fixed time limit. The present European search report has been drawn up for those parts of the European patent application which relate to the invention first mentioned in the claims, namely claims:





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**LACK OF UNITY OF INVENTION  
SHEET B**

Application Number  
**EP 98 10 5398**

The Search Division considers that the present European patent application does not comply with the requirements of unity of invention and relates to several inventions or groups of inventions, namely:

**1. Claims: 1-3, 6, 7**

System for assign a transmission priority to two devices transmitting on a transmission medium, includes determining the position of a collision point on the transmission medium.

**2. Claim : 4**

System for reducing the number of collisions on a transmission medium, including making each transmitter wait before retransmission for one of two different time intervals, one of which is adjustable.

**3. Claim : 5**

Transmission method for a system including repeaters, includes preventing a device from transmitting when a special signal is asserted by a repeater on a local line.

**ANNEX TO THE EUROPEAN SEARCH REPORT  
ON EUROPEAN PATENT APPLICATION NO.**

EP 98 10 5398

This annex lists the patent family members relating to the patent documents cited in the above-mentioned European search report. The members are as contained in the European Patent Office EDP file on  
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26-03-2003

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For more details about this annex : see Official Journal of the European Patent Office, No. 12/82